The uniqueness problem for chromatic symmetric functions of trees

Jeremy L. Martin (University of Kansas)

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## Colorings and the Chromatic Polynomial

Throughout, G = (V, E) will be a simple graph with |V| = n.

A coloring of G is a function  $f: V \to \mathbb{N}$  such that

$$vw \in E \implies f(v) \neq f(w).$$

The chromatic function is  $\chi_G(k) = \#$  colorings  $V \to \{1, \ldots, k\}$ .

Well-known facts:

- $\chi_G(k)$  is a polynomial in k.
- If G is a tree then  $\chi_G(k) = k(k-1)^{n-1}$ .

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- $\chi_G(k)$  is a polynomial in k.
- If G is a tree then  $\chi_G(k) = k(k-1)^{n-1}$ .
- Therefore, the chromatic polynomial of a tree contains no information about it other than the number of vertices.

#### Definition

A symmetric function is a formal power series  $F \in \mathbb{Q}[[x_1, x_2, ...]]$  that is invariant with respect to all permutations of the variables.

Definition Let  $\lambda = (\lambda_1 \dots \lambda_\ell) \vdash n$ . The monomial symmetric function is

 $m_{\lambda} = \text{sum of all monomials of the form } x_{i_1}^{\lambda_1} \cdots x_{i_\ell}^{\lambda^\ell}.$ 

#### Example

$$m_n = \sum_{i=1}^{\infty} x_i^n \qquad m_{11\cdots 1} = \sum_{i_1 < \cdots < i_n} x_{i_1} \cdots x_{i_n}$$

## The Chromatic Symmetric Function

Let f be a coloring of G = (V, E).

Record the number of times each color is used by a monomial:

$$\mathbf{x}^f \stackrel{\text{def}}{=} \prod_{v \in V} x_{f(v)}.$$

The chromatic symmetric function of G is the formal power series

$$X(G) = X_G(x_1, x_2, \dots) = \sum_{\text{colorings } f} \mathbf{x}^f.$$

It was introduced by Stanley in 1995.

# The Chromatic Symmetric Function

- ► X(G) is a well-defined formal power series, because there are only finitely many colorings with a given "palette".
- ► X(G) is a symmetric function because permuting the colors does not change whether a coloring is proper.
- X(G) is homogeneous of degree n = |V(G)|.
- ► The chromatic function \(\chi\_G(k)\) can be recovered from \(X(G): set \(x\_1 = \cdots = x\_k = 1\), and \(x\_i = 0\) for all \(i > k\).

# Examples

- 1.  $X(K_n) = n! \cdot m_{11\dots 1} = \text{sum of all squarefree monomials}$ (= elementary symmetric function  $e_n$ )
- 2.  $X(\overline{K_n}) = m_1^n$ . In general X(G + H) = X(G)X(H).
- 3. These graphs have the same chromatic symmetric function:



4. These two don't:



# Chromatic Symmetric Functions of Trees

For the two trees on 4 vertices...



$$X(P_4) = 24m_{1111} + 6m_{211} + 2m_{22}$$
  
$$X(S_4) = 24m_{1111} + 6m_{211} + m_{31}$$

Question (Stanley)

Do there exist two non-isomorphic trees with the same CSF?

#### Power-Sum Symmetric Functions

The power-sum symmetric functions are

$$p_n = \sum_{i=1}^{\infty} x_i^n, \qquad p_{(\lambda_1,\ldots,\lambda_\ell)} = \prod_{j=1}^{\ell} p_{\lambda_j}.$$

**Fact:** The set  $\{p_{\lambda} : \lambda \vdash n\}$  is a vector space basis for **Sym**<sub>*n*</sub>.

For  $S \subseteq E(G)$ , let type(S) = partition whose parts are the sizes of the connected components of the subgraph  $(V, S) \subseteq G$ .



#### Power-Sum Coefficients of the CSF

Theorem (Stanley '95)

$$X_G = \sum_{S \subseteq E} (-1)^{|S|} p_{\operatorname{type}(S)}.$$

▶ There is cancellation in Stanley's formula iff *G* has cycles.

Corollary If G is a tree and X(G) is written in the basis  $\{p_{\lambda}\}$ , then

coefficient of 
$$p_{\lambda} = (-1)^{n-\ell(\lambda)} \# \{A \subseteq E : \operatorname{type}(A) = \lambda\}.$$

#### Definition The subtree polynomial of a tree T is

$$S_T(q,r) = \sum_U q^{|E(U)|} r^{|L(U)|} = \sum_{i,j} \sigma_{i,j}(T) q^i r^j$$

where U ranges over all subtrees of T with at least one edge, and L(U) is the set of leaf edges.

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Theorem (JLM–Matthew Morin–Jennifer Wagner '06) The numbers  $\sigma_{i,j}(T)$  are linear combinations of the  $c_{\lambda}(T)$ , with coefficients that depend only on n (not on T).

In particular, X(T) determines  $S_T(q, r)$ .

#### Corollary

The degree sequence and distance sequence of T are determined by its chromatic symmetric function.

Proof sketch. Observe that

> $\sigma_{i,i}$  = number of *i*-edge stars in *T*  $\sigma_{i,2}$  = number of *i*-edge paths in *T*

Now use inclusion/exclusion to obtain number of vertices with degree i, and number of pairs of vertices at distance i.

This result is enough to prove that some very special classes of trees (e.g., spiders) are determined by their CSFs.

The smallest non-isomorphic trees with the same subtree polynomial have 11 vertices:



#### Theorem (Keeler Russell, 2013)

Every tree with  $n \le 25$  vertices is determined up to isomorphism by its chromatic symmetric function.

Keeler's proof was entirely a brute-force computation, but with several wrinkles.

- ► How do you generate all ≈ 10<sup>8</sup> trees on 25 vertices? (Hint: Do not use the Prüfer code followed by isomorphism testing.)
- ► Trick 1: Classify trees by degree sequence and parallelize
- Trick 2: Compute and compare one coefficient at a time instead of the entire CSF

The chromatic symmetric function does not obey a deletioncontraction recurrence, but it does satisfy the modular relation:

Theorem (Guay-Paquet '13+; Orellana–Scott '14) Suppose that e, e', e'' form a triangle in G. Then

$$X(G) + X(G - e - e') = X(G - e) + X(G - e').$$

#### Question

Are there other linear relations between the chromatic symmetric functions of trees (or graphs)?

The theory of combinatorial Hopf algebras may be useful...

# Thanks for listening!

# Appendix A: Explicit Formula for the Subtree Polynomial

Subtree polynomial:

$$S_{\mathcal{T}}(q,r) = \sum_{U} q^{|\mathcal{E}(U)|} r^{|\mathcal{L}(U)|} = \sum_{i,j} \sigma_{i,j}(T) q^{i} r^{j}$$

Formula [JLM, Morin, Wagner]:

$$\sigma_{i,j} = \sum_{\lambda \vdash n} (-1)^{i+j} \binom{\ell-1}{\ell-n+i} \sum_{d=1}^{j} \binom{i-d}{j-d} \sum_{k=1}^{\ell} \binom{\lambda_k-1}{d} c_{\lambda}(T)$$

where  $\ell = \text{length of } \lambda$  and  $c_{\lambda}(T) = \text{coefficient of } p_{\lambda} \text{ in } X(T)$ .

(Are you happy you asked?)