# The uniqueness problem for chromatic symmetric functions of trees 

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## Colorings and the Chromatic Polynomial

Throughout, $G=(V, E)$ will be a simple graph with $|V|=n$.
A coloring of $G$ is a function $f: V \rightarrow \mathbb{N}$ such that

$$
v w \in E \quad \Longrightarrow \quad f(v) \neq f(w)
$$

The chromatic function is $\chi_{G}(k)=\#$ colorings $V \rightarrow\{1, \ldots, k\}$.
Well-known facts:

- $\chi_{G}(k)$ is a polynomial in $k$.
- If $G$ is a tree then $\chi_{G}(k)=k(k-1)^{n-1}$.


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- $\chi_{G}(k)$ is a polynomial in $k$.
- If $G$ is a tree then $\chi_{G}(k)=k(k-1)^{n-1}$.
- Therefore, the chromatic polynomial of a tree contains no information about it other than the number of vertices.


## Symmetric Functions

## Definition

A symmetric function is a formal power series $F \in \mathbb{Q}\left[\left[x_{1}, x_{2}, \ldots\right]\right]$ that is invariant with respect to all permutations of the variables.

## Definition

Let $\lambda=\left(\lambda_{1} \ldots \lambda_{\ell}\right) \vdash n$. The monomial symmetric function is

$$
m_{\lambda}=\text { sum of all monomials of the form } x_{i_{1}}^{\lambda_{1}} \cdots x_{i_{\ell}}^{\lambda^{\ell}} \text {. }
$$

Example

$$
m_{n}=\sum_{i=1}^{\infty} x_{i}^{n} \quad m_{11 \cdots 1}=\sum_{i_{1}<\cdots<i_{n}} x_{i_{1}} \cdots x_{i_{n}}
$$

## The Chromatic Symmetric Function

Let $f$ be a coloring of $G=(V, E)$.

Record the number of times each color is used by a monomial:

$$
\mathbf{x}^{f} \stackrel{\text { def }}{=} \prod_{v \in V} x_{f(v)}
$$

The chromatic symmetric function of $G$ is the formal power series

$$
X(G)=X_{G}\left(x_{1}, x_{2}, \ldots\right)=\sum_{\text {colorings } f} \mathbf{x}^{f}
$$

It was introduced by Stanley in 1995.

## The Chromatic Symmetric Function

- $X(G)$ is a well-defined formal power series, because there are only finitely many colorings with a given "palette".
- $X(G)$ is a symmetric function because permuting the colors does not change whether a coloring is proper.
- $X(G)$ is homogeneous of degree $n=|V(G)|$.
- The chromatic function $\chi_{G}(k)$ can be recovered from $X(G)$ : set $x_{1}=\cdots=x_{k}=1$, and $x_{i}=0$ for all $i>k$.


## Examples

1. $X\left(K_{n}\right)=n!\cdot m_{11 \cdots 1}=$ sum of all squarefree monomials (= elementary symmetric function $e_{n}$ )
2. $X\left(\overline{K_{n}}\right)=m_{1}^{n}$. In general $X(G+H)=X(G) X(H)$.
3. These graphs have the same chromatic symmetric function:

4. These two don't:


## Chromatic Symmetric Functions of Trees

For the two trees on 4 vertices...


$$
\begin{aligned}
& X\left(P_{4}\right)=24 m_{1111}+6 m_{211}+2 m_{22} \\
& X\left(S_{4}\right)=24 m_{1111}+6 m_{211}+m_{31}
\end{aligned}
$$

Question (Stanley)
Do there exist two non-isomorphic trees with the same CSF?

## Power-Sum Symmetric Functions

The power-sum symmetric functions are

$$
p_{n}=\sum_{i=1}^{\infty} x_{i}^{n}, \quad p_{\left(\lambda_{1}, \ldots, \lambda_{\ell}\right)}=\prod_{j=1}^{\ell} p_{\lambda_{j}}
$$

Fact: The set $\left\{p_{\lambda}: \lambda \vdash n\right\}$ is a vector space basis for $\mathbf{S y m}_{n}$.
For $S \subseteq E(G)$, let type $(S)=$ partition whose parts are the sizes of the connected components of the subgraph $(V, S) \subseteq G$.

## Example:



$$
\text { type }=(3,3,2,1,1)
$$

## Power-Sum Coefficients of the CSF

Theorem (Stanley '95)

$$
X_{G}=\sum_{S \subseteq E}(-1)^{|S|} p_{\text {type }(S)}
$$

- There is cancellation in Stanley's formula iff $G$ has cycles.

Corollary
If $G$ is a tree and $X(G)$ is written in the basis $\left\{p_{\lambda}\right\}$, then

$$
\text { coefficient of } p_{\lambda}=(-1)^{n-\ell(\lambda)} \#\{A \subseteq E: \text { type }(A)=\lambda\}
$$

## The Subtree Polynomial

## Definition

The subtree polynomial of a tree $T$ is

$$
S_{T}(q, r)=\sum_{U} q^{|E(U)|} r^{|L(U)|}=\sum_{i, j} \sigma_{i, j}(T) q^{i} r^{j}
$$

where $U$ ranges over all subtrees of $T$ with at least one edge, and $L(U)$ is the set of leaf edges.

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Theorem (JLM-Matthew Morin-Jennifer Wagner '06)
The numbers $\sigma_{i, j}(T)$ are linear combinations of the $c_{\lambda}(T)$, with coefficients that depend only on $n$ (not on $T$ ).

In particular, $X(T)$ determines $S_{T}(q, r)$.

## Degree and Distance

## Corollary

The degree sequence and distance sequence of $T$ are determined by its chromatic symmetric function.

Proof sketch.
Observe that

$$
\begin{aligned}
\sigma_{i, i} & =\text { number of } i \text {-edge stars in } T \\
\sigma_{i, 2} & =\text { number of } i \text {-edge paths in } T
\end{aligned}
$$

Now use inclusion/exclusion to obtain number of vertices with degree $i$, and number of pairs of vertices at distance $i$.

## Does This Help Answer Stanley's Question?

This result is enough to prove that some very special classes of trees (e.g., spiders) are determined by their CSFs.

The smallest non-isomorphic trees with the same subtree polynomial have 11 vertices:


## A Heroic Computation

Theorem (Keeler Russell, 2013)
Every tree with $n \leq 25$ vertices is determined up to isomorphism by its chromatic symmetric function.

Keeler's proof was entirely a brute-force computation, but with several wrinkles.

- How do you generate all $\approx 10^{8}$ trees on 25 vertices? (Hint: Do not use the Prüfer code followed by isomorphism testing.)
- Trick 1: Classify trees by degree sequence and parallelize
- Trick 2: Compute and compare one coefficient at a time instead of the entire CSF


## The Modular Relation

The chromatic symmetric function does not obey a deletioncontraction recurrence, but it does satisfy the modular relation:

Theorem (Guay-Paquet '13+; Orellana-Scott '14)
Suppose that $e, e^{\prime}, e^{\prime \prime}$ form a triangle in $G$. Then

$$
X(G)+X\left(G-e-e^{\prime}\right)=X(G-e)+X\left(G-e^{\prime}\right)
$$

## Question

Are there other linear relations between the chromatic symmetric functions of trees (or graphs)?

The theory of combinatorial Hopf algebras may be useful. . .

## Thanks for listening!

## Appendix A: Explicit Formula for the Subtree Polynomial

Subtree polynomial:

$$
S_{T}(q, r)=\sum_{U} q^{|E(U)|} r^{|L(U)|}=\sum_{i, j} \sigma_{i, j}(T) q^{i} r^{j}
$$

Formula [JLM, Morin, Wagner]:

$$
\sigma_{i, j}=\sum_{\lambda \vdash n}(-1)^{i+j}\binom{\ell-1}{\ell-n+i} \sum_{d=1}^{j}\binom{i-d}{j-d} \sum_{k=1}^{\ell}\binom{\lambda_{k}-1}{d} c_{\lambda}(T)
$$

where $\ell=$ length of $\lambda$ and $c_{\lambda}(T)=$ coefficient of $p_{\lambda}$ in $X(T)$.
(Are you happy you asked?)

