A positivity phenomenon in Elser's Gaussian-cluster percolation model

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Definition

A **nucleus** of a connected graph G is a connected subgraph $N \subseteq G$ such that V(N) is a vertex cover.



Definition

Let $\mathcal{N}(G)$ denote the set of all nuclei. The k^{th} Elser number of G is

$$els_k(G) = (-1)^{|V(G)|+1} \sum_{N \in \mathcal{N}(G)} (-1)^{|E(N)|} |V(N)|^k.$$

Example $\operatorname{els}_k(K_2) = 2^k - 2$

Idea: Interpret Elser numbers as Euler characteristics.

Motivation: Percolation Theory

- Percolation theory models a physical medium by an Erdös-Rényi random subgraph Γ of Z^d or some other periodic lattice [H. Kesten, Notices AMS 2006]
- Typically, different values of d have to be studied separately (think the Drunkard's Walk)
- V. Elser [J. Phys. A 1984] proposed a random geometric graph model in which d can be treated as a parameter
- The numbers els_k(G) arise in a generating function for connected components of Γ.
- ► Elser proved that els₁(G) = 0 for all G and conjectured (based on experimentation) that els_k(G) ≥ 0 for all k ≥ 2.

The Main Theorem

Theorem [D-B, H, L, M, N, V, W 2019⁺]

Let G be a connected graph with at least two vertices. Then:

- 1. $els_0(G) \le 0$.
- 2. $els_1(G) = 0$.
- 3. $\operatorname{els}_k(G) \ge 0$ for all $k \ge 2$ (Elser's conjecture).

More specifically:

- 4. If G is 2-connected, then $els_0(G) < 0$ and $els_k(G) > 0$ for all $k \ge 2$.
- 5. Otherwise, $els_k(G) \neq 0$ if and only if $k \geq \ell$, where $\ell \geq 2$ is the number of *leaf blocks* (2-connected components of G that contain exactly one cut-vertex).

Also:

6. (Monotonicity) If $e \in E(G)$ is neither a loop or cut-edge, then

$${\sf els}_k(G) \geq {\sf els}_k(G/e) + {\sf els}_k(G\setminus e)$$

with equality for k = 0.

Nuclei are stable under adding edges. Therefore:

Definition

The **nucleus complex** of G is the simplicial complex on E(G) given by

$$\Delta^{G} = \{ E(\overline{N}) : N \in \mathcal{N}(G) \}.$$

For $U \subseteq V(G)$, the *U***-nucleus complex** is the subcomplex

$$\Delta_U^G = \{ E(\overline{N}) : N \in \mathcal{N}(G), V(N) \supseteq U \}.$$

Annoying Special Case: We have to define $\Delta_{\emptyset}^{cK_2}$ as a non-simplicial Δ -complex, since not all nuclei of cK_2 are determined by their edge sets.

What do the simplicial complexes Δ_U^G look like?

•
$$\Delta_{\emptyset}^{G} = \Delta^{G}$$

• $\Delta_{V(G)}^{G}$ = matroid complex of cographic matroid

▶ In general, *U*-nucleus complexes are not pure or (nonpure) shellable

Proposition

$$\begin{aligned} \mathsf{els}_{k}(G) &= (-1)^{|E(G)| + |V(G)| + 1} \sum_{U \subseteq V(G)} \mathsf{Sur}(k, |U|) \sum_{\substack{N \in \mathcal{N}(G):\\ U \subseteq V(N)}} (-1)^{|E(\overline{N})|} \\ &= (-1)^{|E(G)| + |V(G)|} \sum_{U \subseteq V(G)} \mathsf{Sur}(k, |U|) \, \tilde{\chi}(\Delta_{U}^{\mathsf{G}}) \end{aligned} \tag{1}$$

where $Sur(a, b) = b! \cdot Stir2(a, b) =$ number of surjections $[a] \rightarrow [b]$.

Let Dep(G) be the **deparallelization** of G: identify all edges in the same parallel class. (Note that Dep(G) can have loops.)

Proposition Let G be a graph and $U \subseteq V(G)$.

1. If G has a loop ℓ , then ℓ is a cone point of Δ_U^G , so $\tilde{\chi}(\Delta_U^G) = 0$.

- 2. Let D = Dep(G). Then $\tilde{\chi}(\Delta_U^D) = (-1)^{|E(G)| |E(D)|} \tilde{\chi}(\Delta_U^G)$.
- 3. Suppose $G \neq K_2$ has a cut-edge *e*. Then:
 - ▶ If e is a leaf edge with leaf x and $x \notin U$, then Δ_{U}^{G} is a cone.

• Otherwise,
$$\Delta_U^G = \Delta_{U/e}^{G/e}$$

In terms of Elser numbers:

- 1'. If G has a loop, then $els_k(G) = 0$ for all k.
- 2'. For all G and k, $els_k(G) = els_k(Dep(G))$.

Corollary Let *T* be a tree with $n \ge 3$ vertices. Let *L* be the set of leaf vertices. Then:

$$\tilde{\chi}(\Delta_U^{\mathsf{T}}) = \begin{cases} 1 - |U| & \text{if } T = K_2, \\ 0 & \text{if } T \neq K_2 \text{ and } L \not\subseteq U, \\ -1 & \text{if } T \neq K_2 \text{ and } L \subseteq U. \end{cases}$$
(2)

Corollary For all $k \ge 1$,

$$\mathsf{els}_k(T) = \sum_{i=0}^{n-|L|} \binom{n-|L|}{i} \mathsf{Sur}(k,|L|+i)$$

In particular, Elser's conjecture is true for trees: $els_k(T) \ge 0$ for $k \ge 1$ (and > 0 iff $k \ge |L|$). **Theorem** Let G be a connected graph with $|V(G)| \ge 2$, let $e \in E(G)$ be neither a loop or cut-edge, and let $U \subseteq V(G)$. Then

$$\tilde{\chi}(\Delta_U^G) = \tilde{\chi}(\Delta_{U/e}^{G/e}) - \tilde{\chi}(\Delta_U^{G\setminus e}).$$

Proof sketch: It's a lot like proving a Tutte polynomial identity. Define a bijection $\psi: 2^{E(G)} \to 2^{E(G \setminus e)} \cup 2^{E(G/e)}$ by

$$\psi(A) = \begin{cases} A \setminus e \subseteq E(G \setminus e) & \text{if } e \in A, \\ A \subseteq E(G/e) & \text{if } e \notin A. \end{cases}$$

Then crank out the recurrence, keeping track of vertex sets and treating the case $G = K_2$ separately.

Idea Apply the deletion/contraction recurrence repeatedly until it bottoms out. The result will be an expression

$$\tilde{\chi}(\Delta_U^G) = \sum_{i=1}^s \varepsilon_i \ \tilde{\chi}(\Delta_{U[T_i]}^{T_i})$$
(3)

where

The list T_1, \ldots, T_s is not an invariant of G, but depends on the choices of edge to delete and contract at each stage of the recurrence.

The calculation is recorded by a thing we call a **restricted deletion/contraction tree** (RDCT).



Observation

 ε_i = number of edges deleted = $(|E(G)| - |E(T_i)|) - (|V(G)| - |V(T_i)|)$

Therefore, formula (3) becomes

$$(-1)^{|E(G)|+|V(G)|}\tilde{\chi}(\Delta_U^G) = -\sum_{i=1}^s \tilde{\chi}(\Delta_{U[T_i]}^{T_i}) \quad \begin{cases} \leq 0 & \text{if } |U| = 0 \\ = 0 & \text{if } |U| = 1 \\ \geq 0 & \text{if } |U| \geq 1 \end{cases}$$

and Elser's conjecture follows from the calculation of Elser numbers for trees (2) and the relationship between Elser numbers and $\tilde{\chi}(\Delta_{II}^{G})$ (1).

Question

When are the inequalities ≤ 0 and ≥ 0 strict?

Proposition

Let \mathscr{B} be any RDCT for G, with leaves T_1, \ldots, T_s . Then:

1.
$$els_0(G) = -\#\{i: T_i \cong K_2\}.$$

- 2. For $k \ge 2$, the following are equivalent:
 - els_k(G) > 0;
 - for some $U \subseteq V(G)$ we have $|U| \leq k$ and $\tilde{\chi}(\Delta_U^G) \neq 0$;
 - for some $i \in [s]$ we have $|L(T_i)| \leq k$.

Problem

This result depends on the choice of \mathscr{B} — we would like a description in terms of *G* itself.

Question

Which tree minors show up as leaves of \mathscr{B} for some \mathscr{B} ?

Definition

An ear decomposition of G is a list of subgraphs R_1, \ldots, R_m such that

1.
$$E(G) = E(R_1) \cup \cdots \cup E(R_m);$$

- 2. R_1 is a cycle; and
- 3. for each i > 1, the graph R_i is a path that meets $R_1 \cup \cdots \cup R_{i-1}$ only at its endpoints.



Fact *G* has an ear decomposition if and only if it is 2-connected (i.e., has no cut-vertex).

Theorem [D-B–H–L–M–N–V–W; F. Petrov]

Let G be a 2-connected graph and let $T \subseteq G$ be a spanning tree.

Then G has an ear decomposition $R_1 \cup \cdots \cup R_m$ such that $|E(R_i) \setminus T| = 1$ for every *i*.

(Proof: constructive algorithm)

Corollary Let *G* be 2-connected.

Then *every* tree minor of G can be realized without contracting a cut-edge or deleting a loop, hence appears as a leaf of some RDCT.

Consequently, $els_k(G) > 0$ for all $k \ge 2$.

Proposition

More generally, $els_k(G) \neq 0$ if and only if G has at most k leaf blocks.

Open Question

To what extent do Elser numbers depend on the matroid of G? (They are *not* matroid invariants.)

Conjecture

Let G be a connected graph and $U \subseteq V(G)$. Then the reduced Betti number $\tilde{\beta}_k(\Delta_U^G)$ is nonzero only if

(i)
$$U = \emptyset$$
 and $k = |E(G)| - |V(G)| - 1$, or

(ii) $|U| \ge 2$ and k = |E(G)| - |V(G)|.

- Can be reduced to the case of 2-connected graphs
- True when U is a vertex cover (using Jakob Jonsson's theory of pseudo-independence complexes)
- Verified computationally for $|V(G)| \leq 6$
- What are the Betti numbers??

Open Question

What significance does our result have for Elser's percolation model???

- You for listening
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