

The Traveling Salesman Problem

Definition: A **complete graph** K_N is a graph with N vertices and an edge between every two vertices.

Definition: A **weighted graph** is a graph in which each edge is assigned a *weight* (representing the time, distance, or cost of traversing that edge).

Definition: A **Hamilton circuit** is a circuit that uses every vertex of a graph once.

Definition: The **Traveling Salesman Problem (TSP)** is the problem of finding a **minimum-weight Hamilton circuit** in K_N .

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Example: Willy decides to visit every Australian city important enough to be listed on [this Wikipedia page](#).

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What route minimizes the total distance he has to travel?

I.e., in this weighted K_{16} , which Hamilton circuit has the smallest total weight?

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BIG PROBLEM: There are 16 vertices, so there are

$$(16 - 1)! = 15! = 1,307,674,368,000$$

Hamilton circuits that each need to be checked.

Solving the TSP Without Brute Force

Idea: At each stage in your tour, choose the closest vertex that you have not visited yet.

This is called the **Nearest-Neighbor Algorithm (NNA)**.

[This spreadsheet](#) shows what happens when Willy uses the NNA to construct a Hamilton circuit (with Sydney as the reference vertex).

The Nearest-Neighbor Algorithm

The result: The Nearest-Neighbor algorithm, using Sydney as the reference vertex, yields the Hamilton circuit

SY → CN → ML → HO → AD → AS → UL → BM → KU
→ DA → MI → CS → MK → BR → AL → PE → SY

whose total weight is **21,049 km**.

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But can Willy do better?

The Repetitive Nearest-Neighbor Algorithm

Observation: Willy can use **any city** as the reference vertex!

That is, Willy can execute the Nearest-Neighbor Algorithm sixteen times, using each city once as the reference vertex.

Then, he can pick the Hamilton circuit with the lowest total weight of these sixteen.

This is called the **Repetitive Nearest-Neighbor Algorithm (RNNA)**.

The Repetitive Nearest-Neighbor Algorithm

Ref. vertex	Hamilton circuit	Weight
AD	AD,ML,HO,CN,SY,BR,MK,CS,MI,AS,UL,BM,KU,DA,PE,AL,AD	18543
AL	AL,PE,BM,KU,DA,AS,UL,AD,ML,HO,CN,SY,BR,MK,CS,MI,AL	19795
AS	AS,UL,BM,KU,DA,MI,CS,MK,BR,SY,CN,ML,HO,AD,AL,PE,AS	18459
BR	BR,MK,CS,MI,AS,UL,BM,KU,DA,AD,ML,HO,CN,SY,AL,PE,BR	22113
BM	BM,KU,DA,AS,UL,AD,ML,HO,CN,SY,BR,MK,CS,MI,PE,AL,BM	19148
CS	CS,MK,BR,SY,CN,ML,HO,AD,AS,UL,BM,KU,DA,MI,PE,AL,CS	22936
CN	CN,SY,ML,HO,AD,AS,UL,BM,KU,DA,MI,CS,MK,BR,AL,PE,CN	21149
DA	DA,KU,BM,UL,AS,MI,CS,MK,BR,SY,CN,ML,HO,AD,AL,PE,DA	18543
HO	HO,ML,CN,SY,BR,MK,CS,MI,AS,UL,BM,KU,DA,AD,AL,PE,HO	20141
KU	KU,DA,AS,UL,BM,MI,CS,MK,BR,SY,CN,ML,HO,AD,AL,PE,KU	18785
MK	MK,CS,MI,AS,UL,BM,KU,DA,AD,ML,HO,CN,SY,BR,AL,PE,MK	23255
ML	ML,HO,CN,SY,BR,MK,CS,MI,AS,UL,BM,KU,DA,AD,AL,PE,ML	20141
MI	MI,AS,UL,BM,KU,DA,CS,MK,BR,SY,CN,ML,HO,AD,AL,PE,MI	20877
PE	PE,AL,BM,KU,DA,AS,UL,AD,ML,HO,CN,SY,BR,MK,CS,MI,PE	19148
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UL	UL,AS,MI,CS,MK,BR,SY,CN,ML,HO,AD,BM,KU,DA,PE,AL,UL	20763

(NNA)

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SY	SY,CN,ML,HO,AD,AS,UL,BM,KU,DA,MI,CS,MK,BR,AL,PE,SY	21049	(NNA)
UL	UL,AS,MI,CS,MK,BR,SY,CN,ML,HO,AD,BM,KU,DA,PE,AL,UL	20763	

The Repetitive Nearest-Neighbor Algorithm

Apparently, using Alice Springs (AS) as the reference vertex yields the best Hamilton circuit so far, namely

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Remember: Willy can still start anywhere he wants!
For instance,

SY → CN → ML → HO → AD → AL → PE → AS
→ UL → BM → KU → DA → MI → CS → MK → BR → SY

represents the same Hamilton circuit.

The Repetitive Nearest-Neighbor Algorithm

Randomly chosen Hamilton circuit:	40,680 km
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- ▶ This algorithm is still **efficient**, but . . .

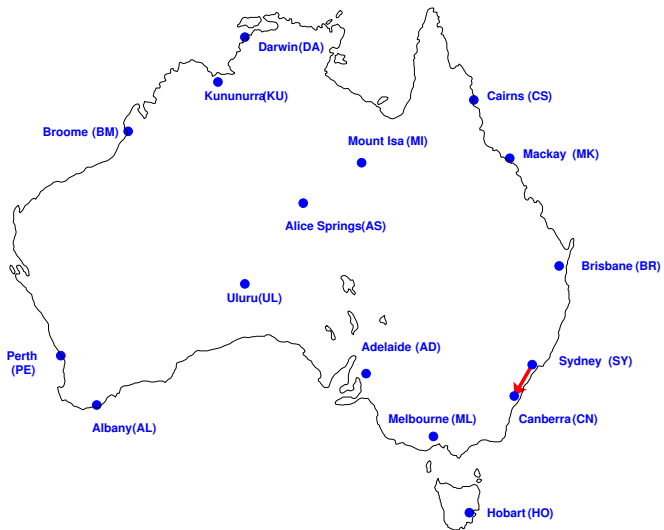
The Repetitive Nearest-Neighbor Algorithm

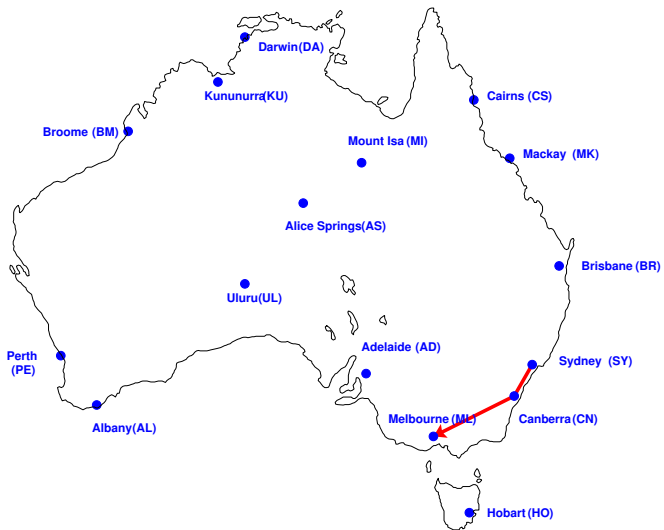
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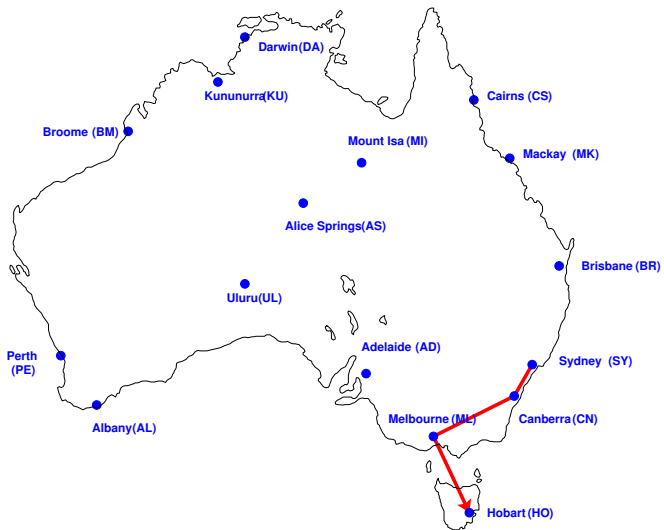
- ▶ In general, there's no way of knowing in advance which reference vertex will yield the best result.
- ▶ This algorithm is still **efficient**, but . . .
- ▶ Is it **optimal**? That is, **Can Willy do even better?**

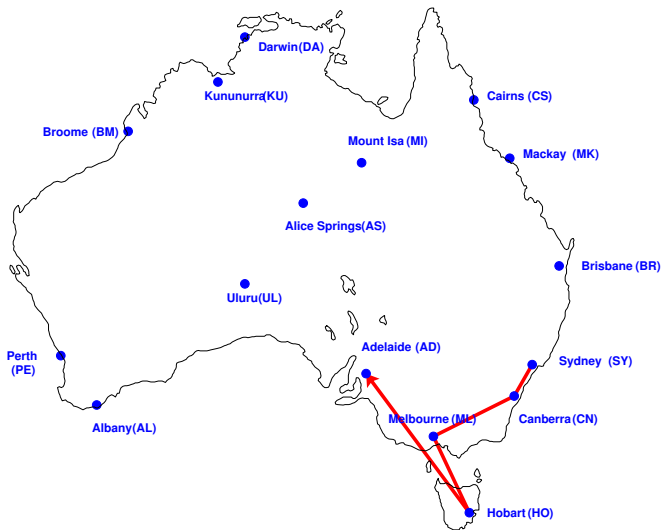
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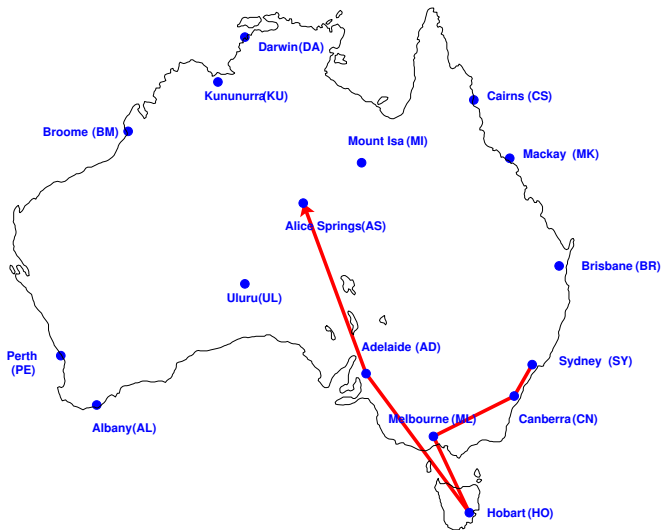
If we look at the map (warning: not quite to scale!) it becomes clear that the RNNA has not produced an optimal Hamilton circuit.

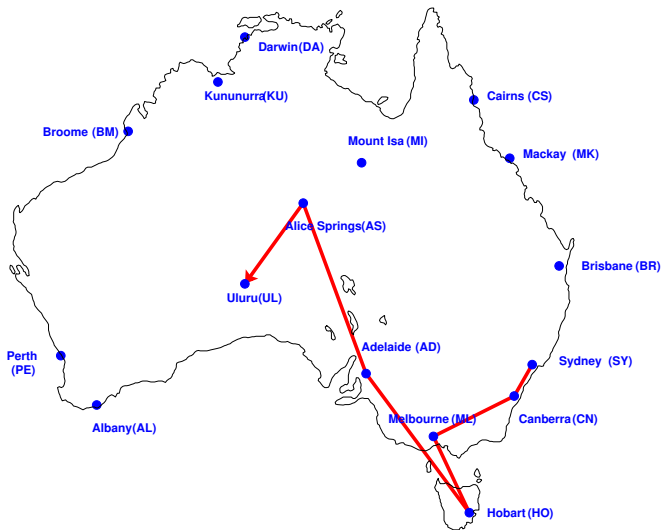


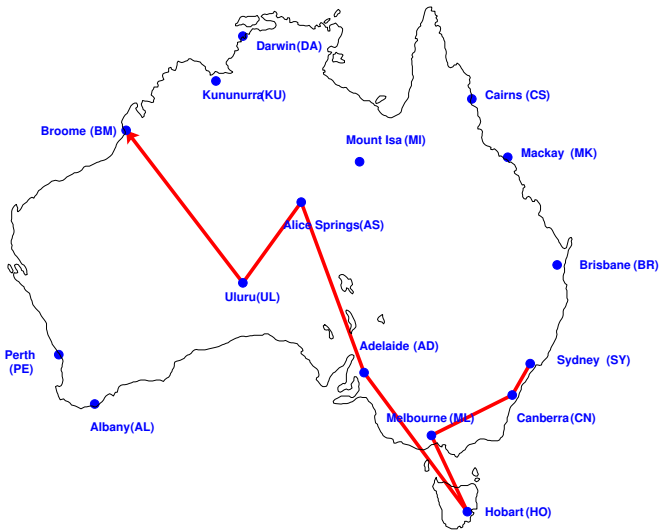




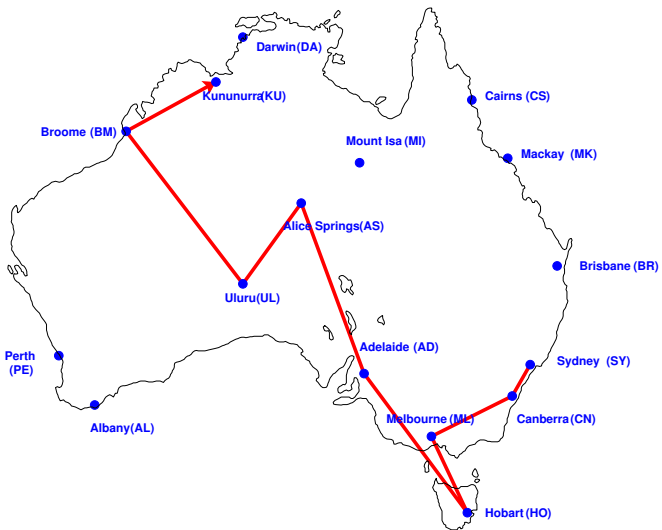


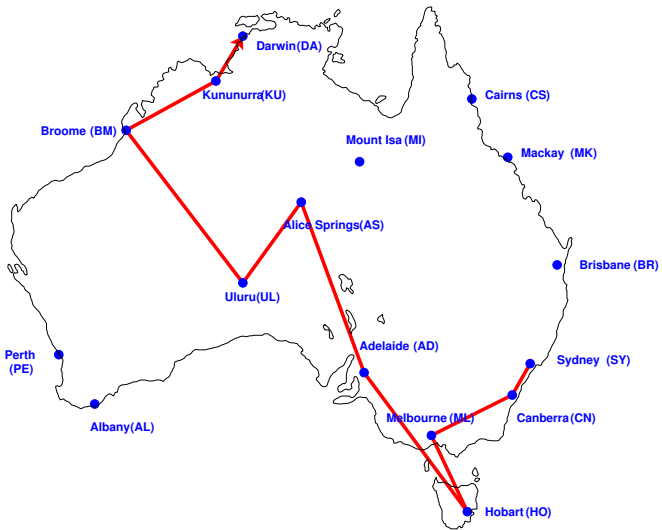


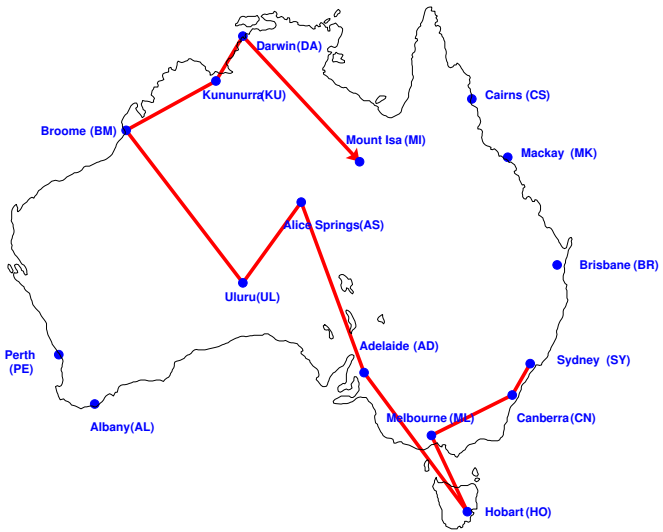


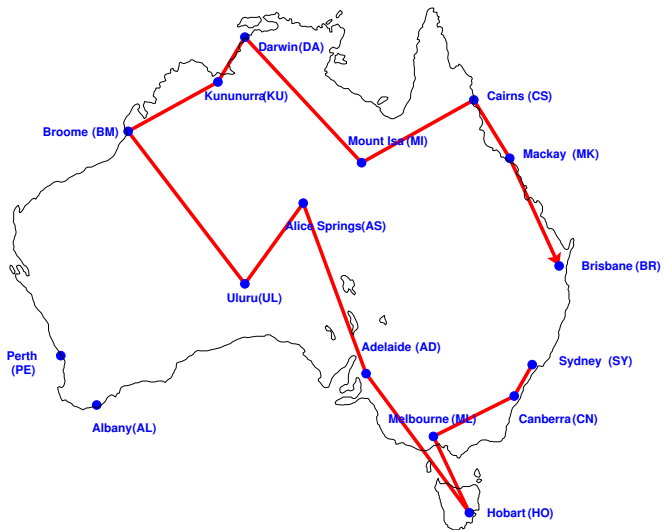


Oops.

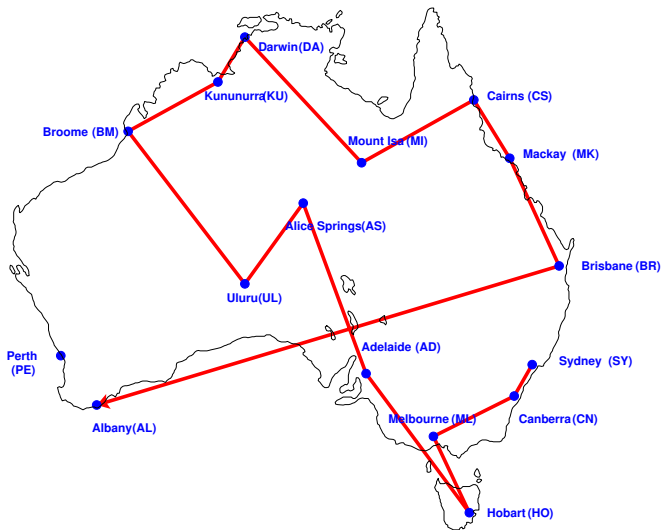


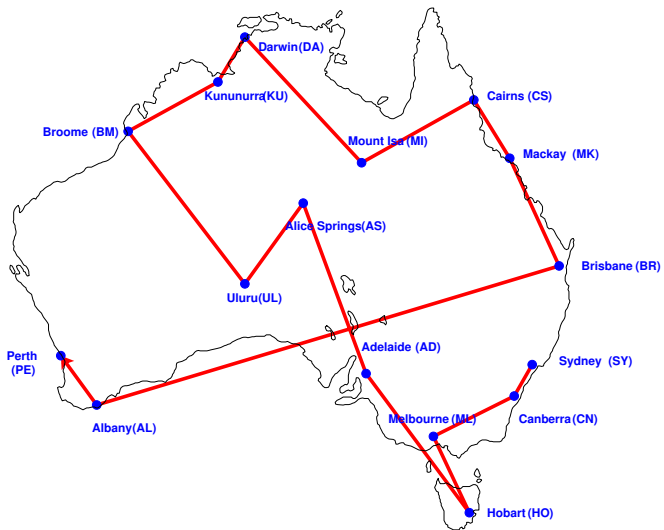


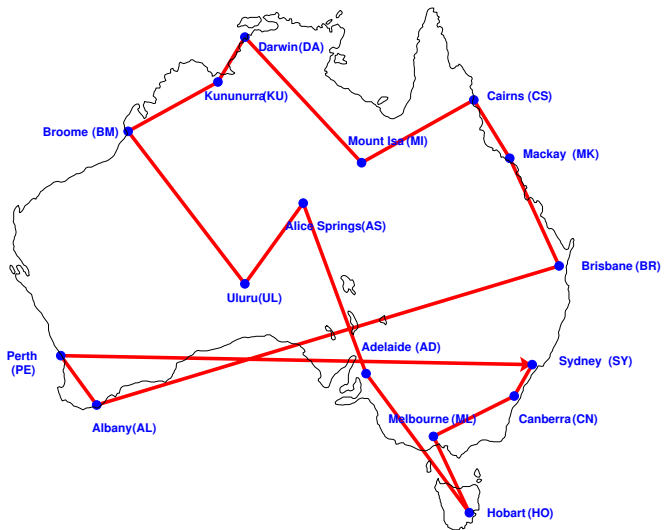




Now the algorithm is stuck — some very expensive edges are required to complete the Hamilton circuit.



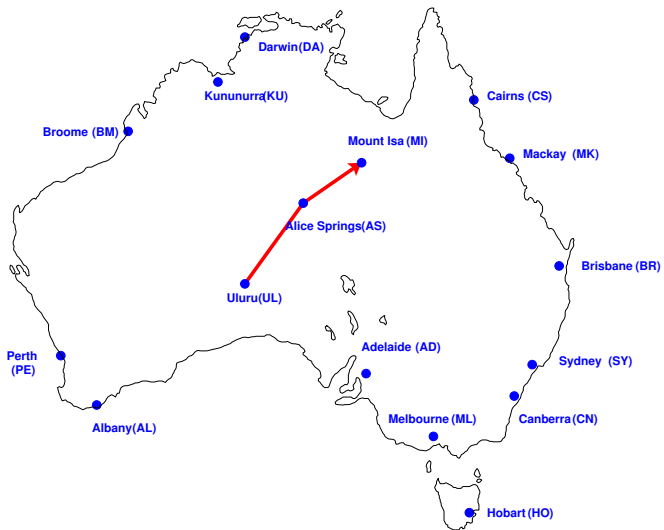




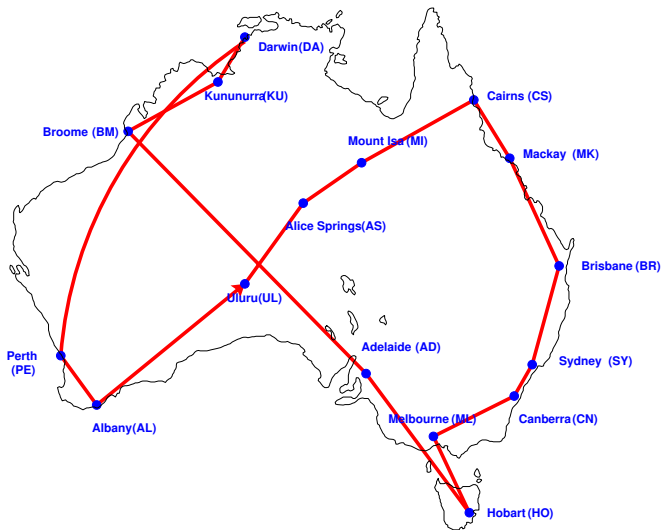
The Repetitive Nearest-Neighbor Algorithm

Starting from Uluru would have created a different problem.









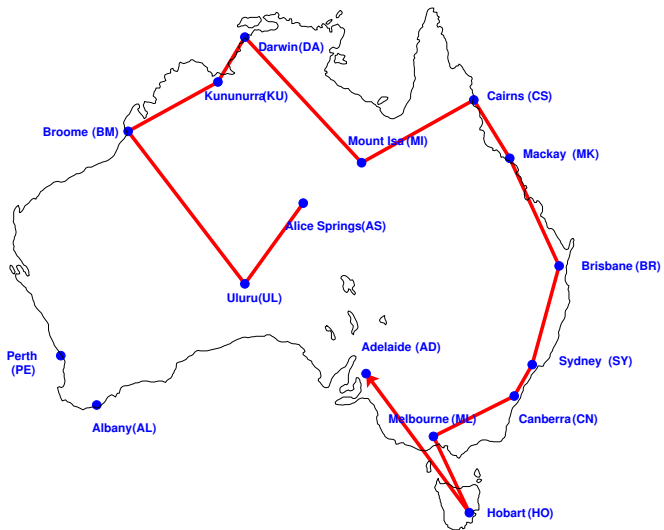
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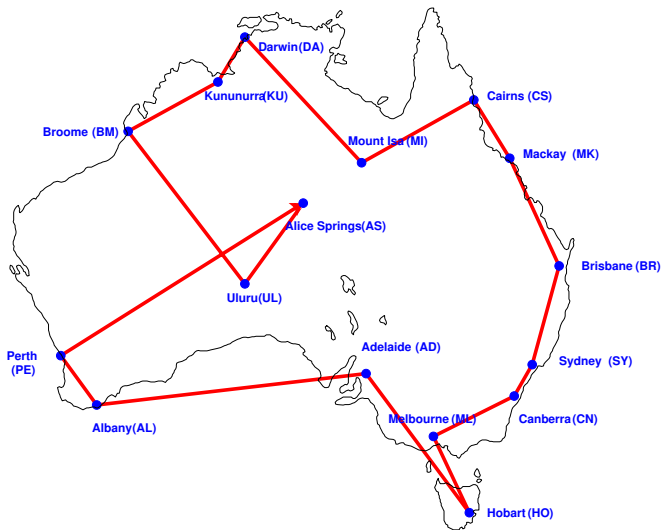
Starting from Alice Springs would have created a different problem (but a less harmful one).



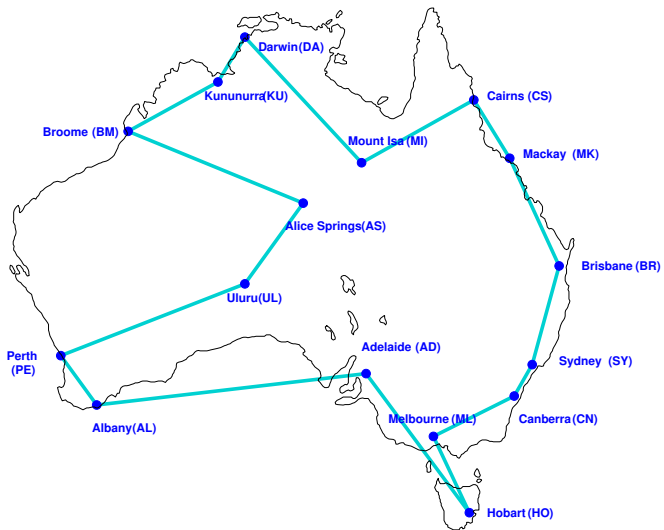








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But how do you make such improvements **part of the algorithm?**

Does starting How about

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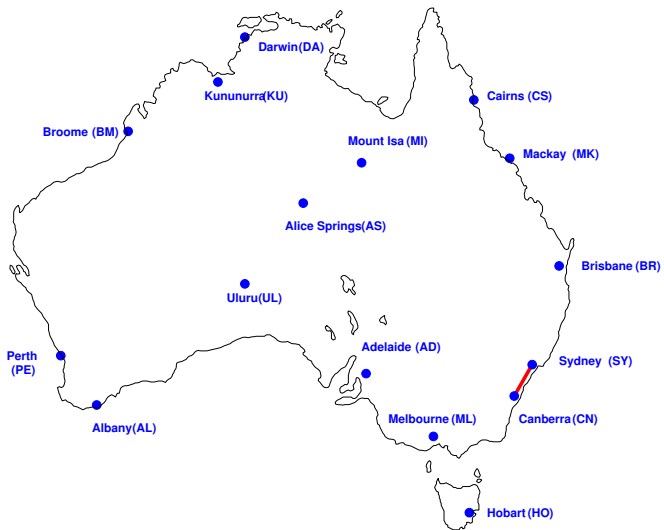
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- ▶ Make sure you add exactly two edges at each vertex.

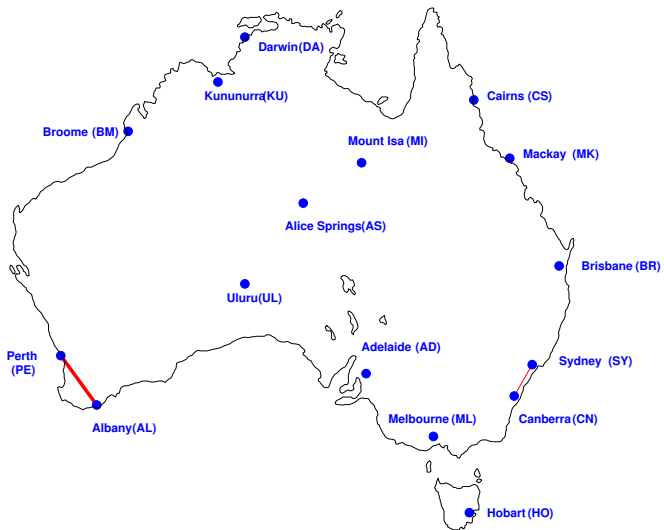
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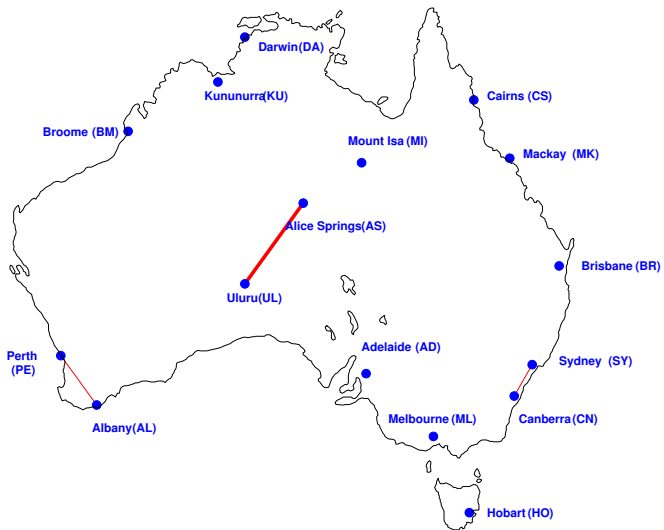
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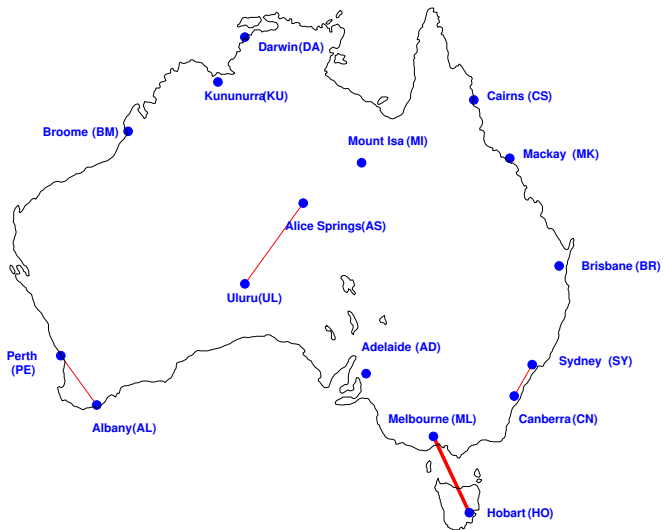
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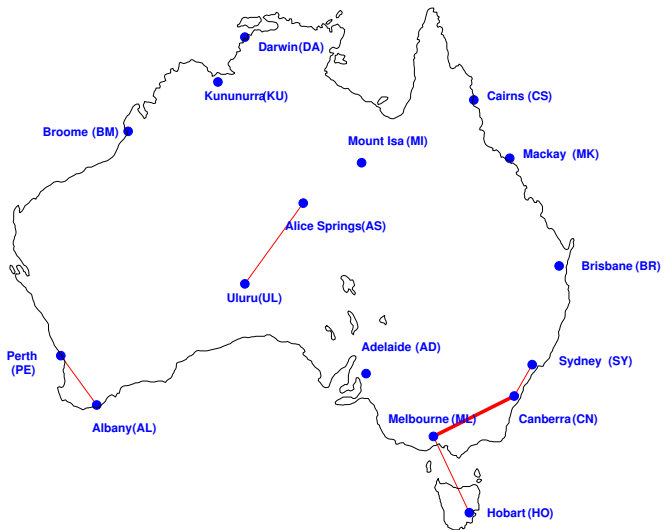
This is called the **Cheapest-Link Algorithm, or CLA.**
[Here is an example.](#)

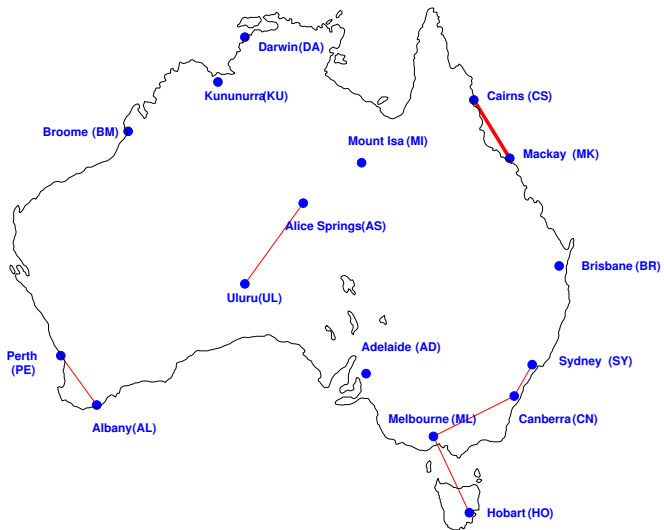


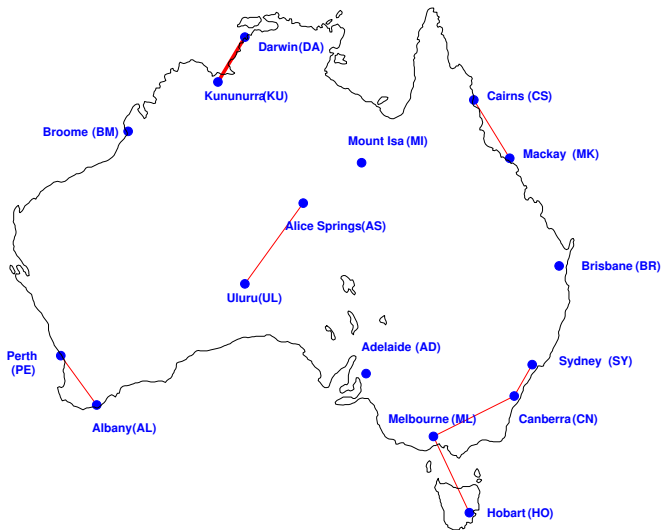


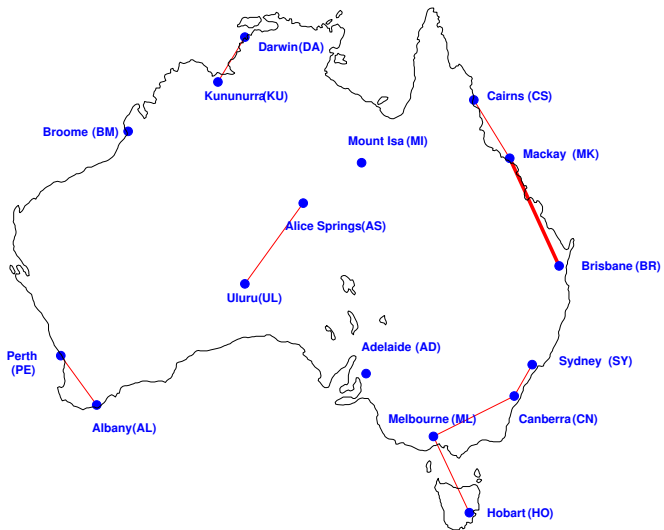


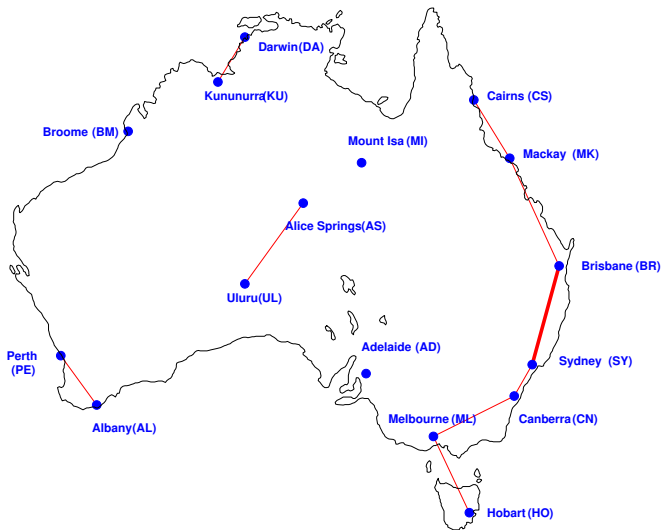


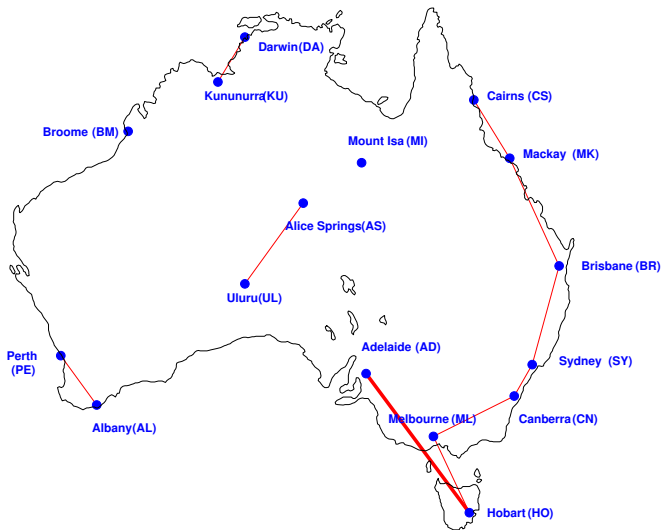


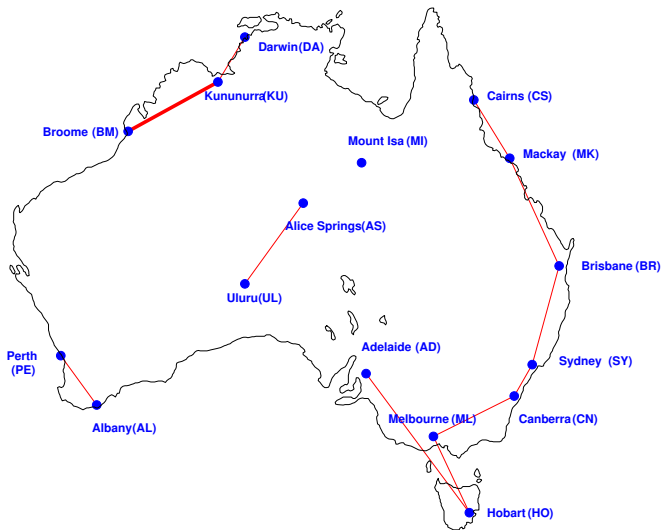


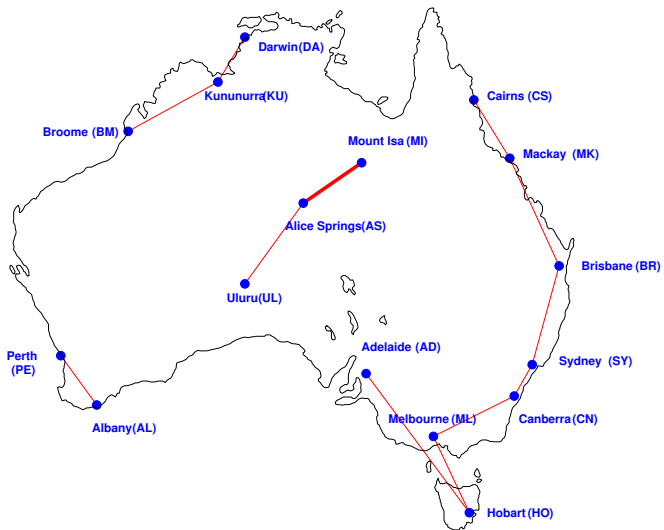


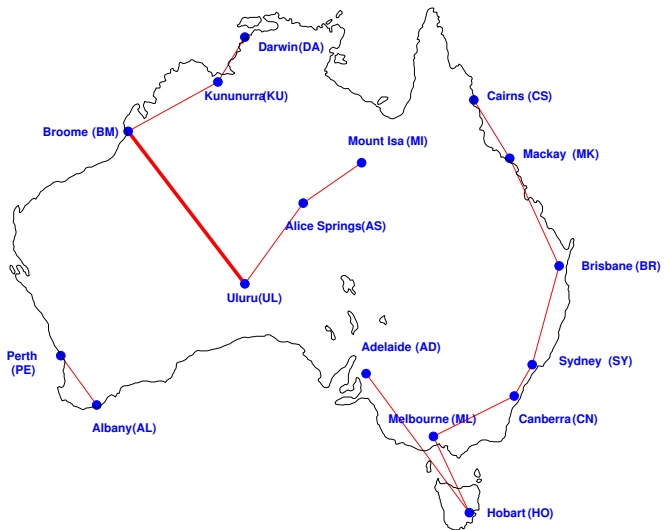


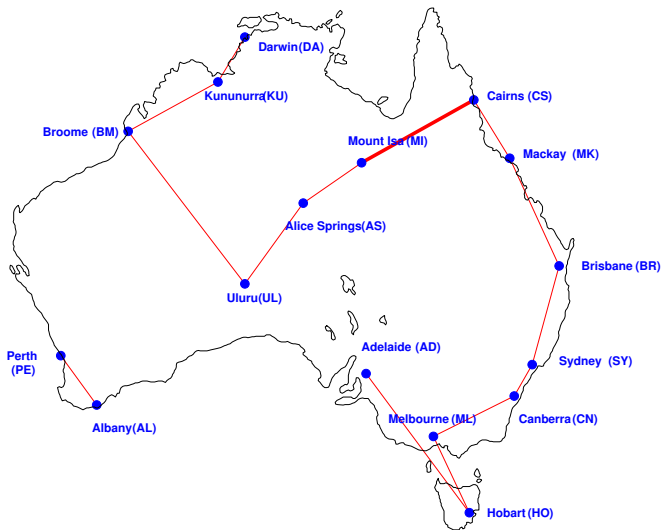


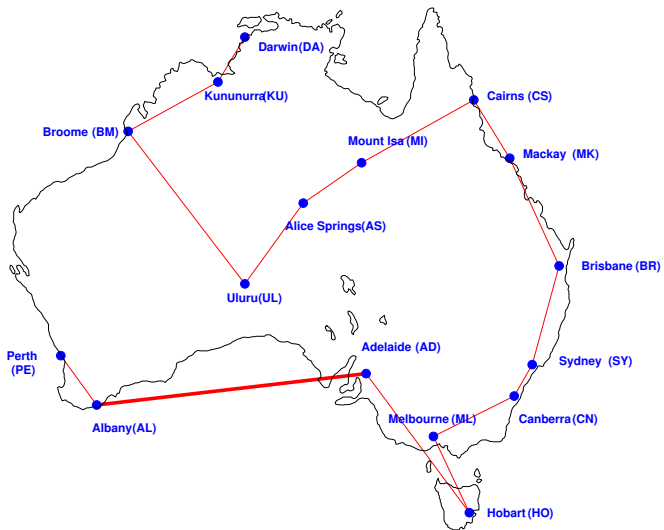


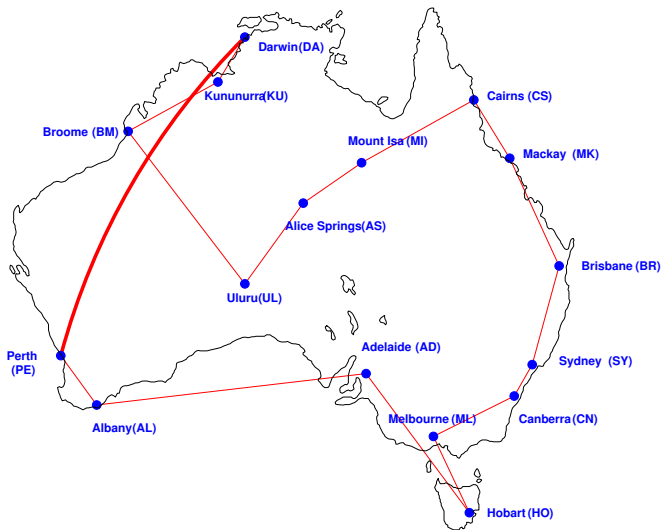












Comparing Algorithms

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- ▶ Brute-force is optimal but not efficient.
- ▶ NNA, RNNA, and CLA are efficient but not optimal.