The Mathematics of Touring (Chapter 6)

- In Chapter 5, we studied Euler paths and Euler circuits: paths and circuits that use every edge of a graph.
- In Chapter 6, we'll look at circuits that use every vertex of a graph exactly once. These are called Hamilton circuits.
- Instead of asking, "Does a graph have a Hamilton circuit?", the interesting question is often, "Out of all the possible Hamilton circuits, which one is the most efficient?"

Willy, a traveling salesman, has to visit each of several cities (say, the 48 state capitals of the continental United States)

He would like his trip to cover as little distance as possible.

In what order should Willy visit the 48 cities?

This problem is the Traveling Salesman Problem, or TSP.

Note that there are many possible routes — to be exact, there are $47! \approx 2.6 \times 10^{59}$ of them. The problem is to find the **shortest** of these routes.

The TSP comes up in many other contexts. For example:

The TSP comes up in many other contexts. For example:

 A spacecraft needs to visit each of six sites on Mars to collect samples (fuel is very expensive on Mars, so shortening the route will save money)

The TSP comes up in many other contexts. For example:

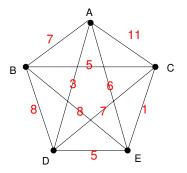
- A spacecraft needs to visit each of six sites on Mars to collect samples (fuel is very expensive on Mars, so shortening the route will save money)
- A school bus needs to visit each of several pickup/dropoff locations (here the issue is not money, but time)

The TSP comes up in many other contexts. For example:

- A spacecraft needs to visit each of six sites on Mars to collect samples (fuel is very expensive on Mars, so shortening the route will save money)
- A school bus needs to visit each of several pickup/dropoff locations (here the issue is not money, but time)
- You are taking your four-year-old trick-or-treating and needs to visit each of eight friends and relatives

The TSP As A Graph Problem

Suppose we have a graph in which every edge has a **weight** (representing its cost, time, or distance).



The TSP is then to find a path or a circuit that

- visits every vertex; and
- has total weight as low as possible.

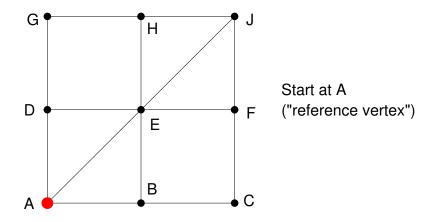
A **Hamilton path** is a path that uses **every vertex** of a graph **exactly once.**

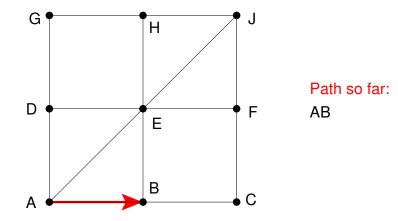
A **Hamilton circuit** is a circuit that uses **every vertex** of a graph **exactly once**.

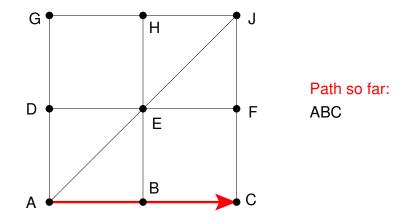
A **Hamilton path** is a path that uses **every vertex** of a graph **exactly once**.

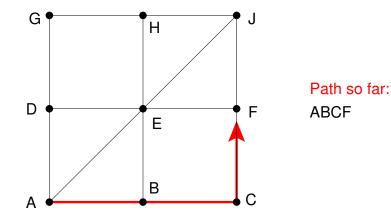
A Hamilton circuit is a circuit that uses every vertex of a graph exactly once.

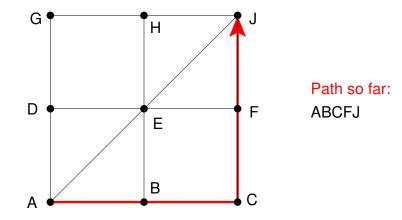
- By contrast, an Euler path/circuit is a path/circuit that uses every *edge* exactly once. (Mnemonic: <u>Euler</u> = <u>E</u>dge)
- "Path": starting and ending vertices are different.
- "circuit" starting and ending vertices are the same.

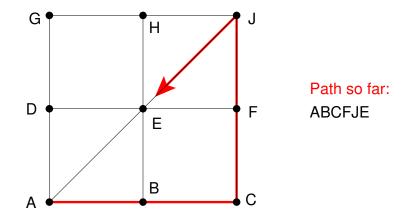


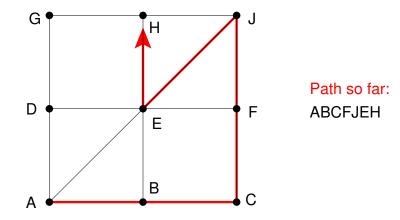


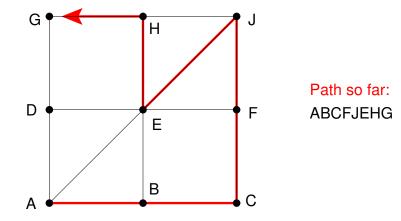


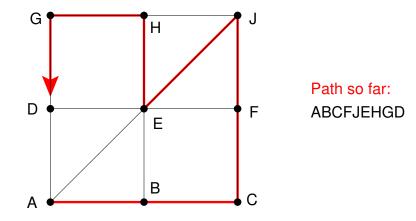


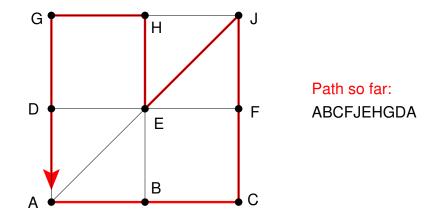


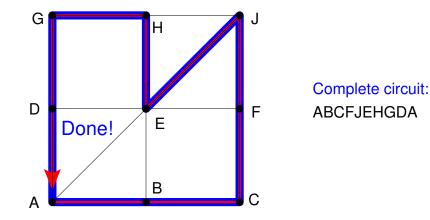




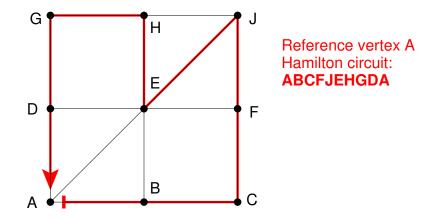


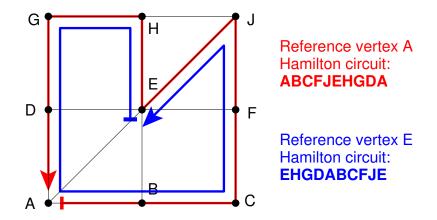




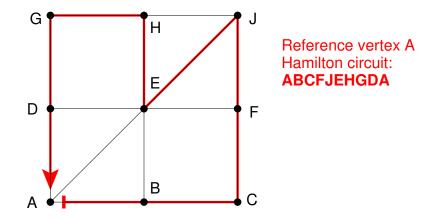


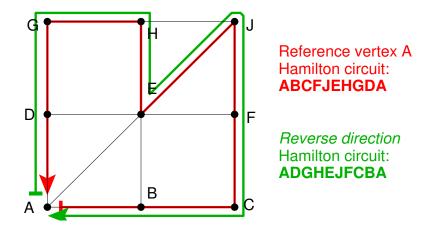
Changing the starting vertex (or "reference vertex") does not change the Hamilton circuit, because the same edges are traversed in the same directions.



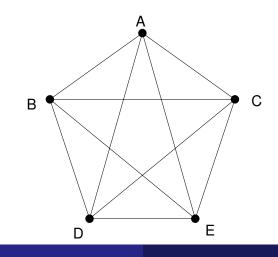


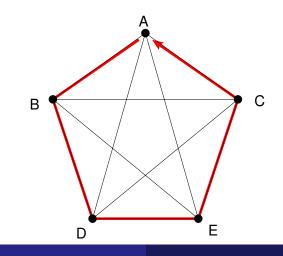
We can also make a Hamilton circuit into its "mirror image" by reversing direction. The mirror image uses the same edges, but **backwards**, so it is not considered the same as the original Hamilton circuit.

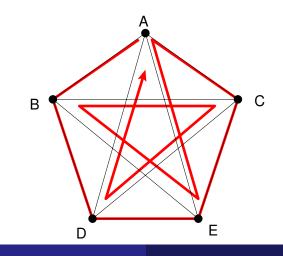








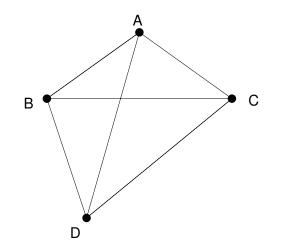






Can a graph have a Hamilton circuit, but not an Euler circuit?

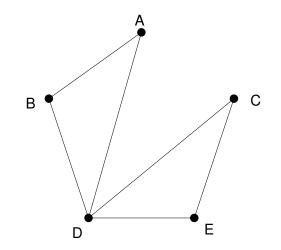
Can a graph have a Hamilton circuit, but not an Euler circuit?





Can a graph have an Euler circuit, but not a Hamilton circuit?

Can a graph have an Euler circuit, but not a Hamilton circuit?



Can a graph have neither a Hamilton circuit nor an Euler circuit? (Let's just consider connected graphs.)

Can a graph have neither a Hamilton circuit nor an Euler circuit? (Let's just consider connected graphs.)



This graph has no circuits at all!

Conclusion: Whether a graph does or does not have a Hamilton circuit **tells you nothing** about whether it has an Euler circuit, and vice versa.

The same is true for Hamilton/Euler **paths** (rather than circuits).

We know how to determine whether a graph has an Euler path or circuit: count the odd vertices.

On the other hand, there is no simple way to tell whether or not a given graph has a Hamilton path or circuit. Rather than asking whether a particular graph has a Hamilton circuit, we will be looking at graphs with **lots** of Hamilton circuits, and trying to find the **shortest** one.

For example, Willy the traveling salesman has the option to drive from any state capital to any other, so the relevant graph has lots of edges — it is a **complete graph**.