## The Traveling Salesman Problem

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Definition: A weighted graph is a graph in which each edge is assigned a weight (representing the time, distance, or cost of traversing that edge).

Definition: The Traveling Salesman Problem is the problem of finding a minimum-weight Hamilton circuit in $K_{N}$.

## The Nearest-Neighbor Algorithm

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3. At each stage in your tour, walk to the nearest neighbor that you have not already visited.
4. When you have visited all vertices, return to the starting vertex.

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- Usually, there is no way to know in advance which reference vertex will work the best.
- Once you find a Hamilton circuit, you can start your tour anywhere you want.


## Example: Willy's Tour of Australia

| Ref. vert | cuit | Weight |
| :---: | :---: | :---: |
|  |  |  |
| AL |  | 19795 |
| AS |  |  |
| BT |  |  |
| BM |  |  |
| CS |  |  |
| CN |  |  |
| DA |  |  |
| но |  |  |
| kU | ku.Da, |  |
| MK |  | 2325 |
| ML |  |  |
| MI | v.oa |  |
| PE |  |  |
|  |  |  |
|  |  |  |

## Example: Willy's Tour of Australia

| Ref. vertex | Hamilton circuit | Weight |
| :---: | :--- | :--- |
| AD | AD,ML,HO,CN,SY,BT,MK,CS,MI,AS,UL,BM,KU,DA,PE,AL,AD | 18543 |
| AL | AL,PE,BM,KU,DA,AS,UL,AD,ML,HO,CN,SY,BT,MK,CS,MI,AL | 19795 |
| AS | AS,UL,BM,KU,DA,MI,CS,MK,BT,SY,CN,ML,HO,AD,AL,PE,AS | 18459 |
| BT | BT,MK,CS,MI,AS,UL,BM,KU,DA,AD,ML,HO,CN,SY,AL,PE,BT | 22113 |
| BM | BM,KU,DA,AS,UL,AD,ML,HO,CN,SY,BT,MK,CS,MI,PE,AL,BM | 19148 |
| CS | CS,MK,BT,SY,CN,ML,HO,AD,AS,UL,BM,KU,DA,MI,PE,AL,CS | 22936 |
| CN | CN,SY,ML,HO,AD,AS,UL,BM,KU,DA,MI,CS,MK,BT,AL,PE,CN | 21149 |
| DA | DA,KU,BM,UL,AS,MI,CS,MK,BT,SY,CN,ML,HO,AD,AL,PE,DA | 18543 |
| HO | HO,ML,CN,SY,BT,MK,CS,MI,AS,UL,BM,KU,DA,AD,AL,PE,HO | 20141 |
| KU | KU,DA,AS,UL,BM,MI,CS,MK,BT,SY,CN,ML,HO,AD,AL,PE,KU | 18785 |
| MK | MK,CS,MI,AS,UL,BM,KU,DA,AD,ML,HO,CN,SY,BT,AL,PE,MK | 23255 |
| ML | ML,HO,CN,SY,BT,MK,CS,MI,AS,UL,BM,KU,DA,AD,AL,PE,ML | 20141 |
| MI | MI,AS,UL,BM,KU,DA,CS,MK,BT,SY,CN,ML,HO,AD,AL,PE,MI | 20877 |
| PE | $P E, A L, B M, K U, D A, A S, U L, A D, M L, H O, C N, S Y, B T, M K, C S, M I, P E ~$ | 19148 |
| SY | SY,CN,ML,HO,AD,AS,UL,BM,KU,DA,MI,CS,MK,BT,AL,PE,SY | 21049 |
| UL | UL,AS,MI,CS,MK,BT,SY,CN,ML,HO,AD,BM,KU,DA,PE,AL,UL | 20763 |

## The Repetitive Nearest-Neighbor Algorithm

Using Alice Springs (AS) as the reference vertex yields the best result:
$\mathrm{AS} \rightarrow \mathrm{UL} \rightarrow \mathrm{BM} \rightarrow \mathrm{KU} \rightarrow \mathrm{DA} \rightarrow \mathrm{MI} \rightarrow \mathrm{CS} \rightarrow \mathrm{MK} \rightarrow \mathrm{BT}$
$\rightarrow \mathrm{SY} \rightarrow \mathrm{CN} \rightarrow \mathrm{ML} \rightarrow \mathrm{HO} \rightarrow \mathrm{AD} \rightarrow \mathrm{AL} \rightarrow \mathrm{PE} \rightarrow \mathrm{AS}$

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$\rightarrow \mathrm{SY} \rightarrow \mathrm{CN} \rightarrow \mathrm{ML} \rightarrow \mathrm{HO} \rightarrow \mathrm{AD} \rightarrow \mathrm{AL} \rightarrow \mathrm{PE} \rightarrow \mathrm{AS}$

Remember: Willy can still start anywhere he wants!
For instance,

$$
\mathrm{SY} \rightarrow \mathrm{CN} \rightarrow \mathrm{ML} \rightarrow \mathrm{HO} \rightarrow \mathrm{AD} \rightarrow \mathrm{AL} \rightarrow \mathrm{PE} \rightarrow \mathrm{AS}
$$

$\rightarrow \mathrm{UL} \rightarrow \mathrm{BM} \rightarrow \mathrm{KU} \rightarrow \mathrm{DA} \rightarrow \mathrm{MI} \rightarrow \mathrm{CS} \rightarrow \mathrm{MK} \rightarrow \mathrm{BT} \rightarrow \mathrm{SY}$
represents the same Hamilton circuit.

## The Cheapest-Link Algorithm

Idea: Start in the middle.

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- Make sure you add exactly two edges at each vertex.
- Don't close the circuit until all vertices are in it.


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- Add the cheapest available edge to your tour. (If there is a tie, break it randomly.)
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- Don't close the circuit until all vertices are in it.

This is called the Cheapest-Link Algorithm, or CLA. Here is an example.

## Example 1



## Results of Example 1

- Output of RNNA: BEDCAB (weight 34)
- Output of CLA: ACBEDA (weight 38)
- In this example, RNNA produces a better result.


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- Output of CLA: ACBEDA (weight 38)
- In this example, RNNA produces a better result.
- In fact, neither of these Hamilton circuits is optimal - the optimal one is EACBDE (weight 32).

Example 2


## Results of Example 2

- RNNA and CLA both output DAECBD (weight 46)
- This happens to be an optimal Hamilton circuit.

Example 3


## Results of Example 3

- Here, the output of both the CLA and the RNNA may depend on how you break ties. (There's no way to know in advance.)

Example 4


Distance table for Example 4

|  | $\mathbf{A}$ | $\mathbf{B}$ | $\mathbf{C}$ | $\mathbf{D}$ | $\mathbf{E}$ | $\mathbf{F}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\mathbf{A}$ |  | 12 | 29 | 22 | 13 | 24 |
| $\mathbf{B}$ | 12 |  | 19 | 3 | 25 | 6 |
| $\mathbf{C}$ | 29 | 19 |  | 21 | 23 | 28 |
| $\mathbf{D}$ | 22 | 3 | 21 |  | 4 | 5 |
| $\mathbf{E}$ | 13 | 25 | 23 | 4 |  | 16 |
| $\mathbf{F}$ | 24 | 6 | 28 | 5 | 16 |  |

## Results of Example 4

- Output of RNNA: FDBAECF (weight 84)
- Output of CLA: ACFBDEA (weight 83)
- In this example, CLA produces a better result.


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- Output of RNNA: FDBAECF (weight 84)
- Output of CLA: ACFBDEA (weight 83)
- In this example, CLA produces a better result.
- Neither of these Hamilton circuits is optimal - the optimal one is FBCAEDF (weight 76).


## Example 5



## Results of Example 5

| Algorithm | Output | Weight |
| :---: | :---: | :---: |
| NNA (A) | ABCDA | $12+15+29+17=73$ |
| NNA (B) | BACDB | $12+14+29+18=73$ |
| NNA (C) | CABDC | $=$ BACDB |
| NNA (D) | DABCD | $=$ ABCDA |
| CLA | ABCDA |  |

## Results of Example 5

| Algorithm | Output | Weight |
| :---: | :---: | :---: |
| NNA (A) | ABCDA | $12+15+29+17=73$ |
| NNA (B) | BACDB | $12+14+29+18=73$ |
| NNA (C) | CABDC | $=$ BACDB |
| NNA (D) | DABCD | $=$ ABCDA |
| CLA | ABCDA |  |

- The only other Hamilton circuit in $K_{4}$ is ACBDA, which has weight $14+15+18+17=64$.


## Results of Example 5

| Algorithm | Output | Weight |
| :---: | :---: | :---: |
| NNA (A) | ABCDA | $12+15+29+17=73$ |
| NNA (B) | BACDB | $12+14+29+18=73$ |
| NNA (C) | CABDC | $=$ BACDB |
| NNA (D) | DABCD | $=$ ABCDA |
| CLA | ABCDA |  |

- The only other Hamilton circuit in $K_{4}$ is ACBDA, which has weight $14+15+18+17=64$.
- So both RNNA and CLA give worst possible answers!


## The Bad News

There is no known algorithm to solve the TSP that is both optimal and efficient.

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There is no known algorithm to solve the TSP that is both optimal and efficient.

- Brute-force is optimal but not efficient.
- NNA, RNNA, and CLA are all efficient but not optimal.

